

CP/M DYNAMIC DEBUGGING TOOL (DDT)

USER'S GUIDE

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CHAPTER I

1870

The first of the year was a very dry one, and the crops were much injured. The weather was very hot, and the ground was very dry. The crops were much injured, and the people were very poor. The first of the year was a very dry one, and the crops were much injured. The weather was very hot, and the ground was very dry. The crops were much injured, and the people were very poor.



# CP/M Dynamic Debugging Tool (DDT)

## User's Guide

### I. Introduction

The DDT program allows dynamic interactive testing and debugging of programs generated in the CP/M environment. The debugger is initiated by typing one of the following commands at the CP/M Console Command level:

```
DDT
DDT filename.HEX
DDT filename.COM
```

where "filename" is the name of the program to be loaded and tested. In both cases, the DDT program is brought into main memory in the place of the Console Command Processor (refer to the CP/M Interface Guide for standard memory organization), and thus resides directly below the Basic Disk Operating System portion of CP/M. The BDOS starting address, which is located in the address field of the JMP instruction at location 5H, is altered to reflect the reduced Transient Program Area size.

The second and third forms of the DDT command shown above perform the same actions as the first, except there is a subsequent automatic load of the specified HEX or COM file. The action is identical to the sequence of commands

```
DDT
Ifilename.HEX or Ifilename.COM
R
```

where the I and R commands set up and read the specified program to test (see the explanation of the I and R commands below for exact details).

Upon initiation, DDT prints the following sign-on message.

```
DDT VERS n.n
```

where nn represents the version number.



		the source file
p2:	A,B, ..., Y	designates the disk name which will receive the hex file
	Z	skips the generation of the hex file
p3:	A,B, ..., Y	designates the disk name which will receive the print file
	X	places the listing at the console
	Z	skips generation of the print file

Thus, the command

ASM X.AAA

indicates that the source file (X.ASM) is to be taken from disk A, and that the hex (X.HEX) and print (X.PRN) files are to be created also on disk A. This form of the command is implied if the assembler is run from disk A. That is, given that the operator is currently addressing disk A, the above command is equivalent to

ASM X

The command

ASM X.ABX

indicates that the source file is to be taken from disk A, the hex file is placed on disk B, and the listing file is to be sent to the console. The command

ASM X.BZZ

takes the source file from disk B, and skips the generation of the hex and print files (this command is useful for fast execution of the assembler to check program syntax).

The source program format is compatible with both the Intel 8080 assembler (macros are not currently implemented in the CP/M assembler, however), as well as the Processor Technology Software Package #1 assembler. That is, the CP/M assembler accepts source programs written in either format. There are certain extensions in the CP/M assembler which make it somewhat easier to use. These extensions are described below.

## 2. PROGRAM FORMAT.

An assembly language program acceptable as input to the assembler consists of a sequence of statements of the form

line#    label    operation    operand    ;comment

where any or all of the fields may be present in a particular instance. Each



## DDT X.COM

which reloads previously saved program from location 100H through page 18 (12FFH). The machine state is not a part of the COM file, and thus the program must be restarted from the beginning in order to properly test it.

## II. DDT COMMANDS.

The individual commands are given below in some detail. In each case, the operator must wait for the prompt character (-) before entering the command. If control is passed to a program under test, and the program has not reached a breakpoint, control can be returned to DDT by executing a RST 7 from the front panel (note that the rubout key should be used instead if the program is executing a T or U command). In the explanation of each command, the command letter is shown in some cases with numbers separated by commas, where the numbers are represented by lower case letters. These numbers are always assumed to be in a hexadecimal radix, and from one to four digits in length (longer numbers will be automatically truncated on the right).

Many of the commands operate upon a "CPU state" which corresponds to the program under test. The CPU state holds the registers of the program being debugged, and initially contains zeroes for all registers and flags except for the program counter (P) and stack pointer (S), which default to 100H. The program counter is subsequently set to the starting address given in the last record of a HEX file if a file of this form is loaded (see the I and R commands).

1. The A (Assemble) Command. DDT allows inline assembly language to be inserted into the current memory image using the A command which takes the form

As

where s is the hexadecimal starting address for the inline assembly. DDT prompts the console with the address of the next instruction to fill, and reads the console, looking for assembly language mnemonics (see the Intel 8080 Assembly Language Reference Card for a list of mnemonics), followed by register references and operands in absolute hexadecimal form. Each successive load address is printed before reading the console. The A command terminates when the first empty line is input from the console.

Upon completion of assembly language input, the operator can review the memory segment using the DDT disassembler (see the L command).

Note that the assembler/disassembler portion of DDT can be overlaid by the transient program being tested, in which case the DDT program responds with an error condition when the A and L commands are used (refer to Section IV).



Technology assembler has the side effect in its operation of ignoring the characters after the operand field has been scanned. This causes an ambiguous situation when attempting to be compatible with Intel's language, since arbitrary expressions are allowed in this case. Hence, programs which use this side effect to introduce comments, must be edited to place a ";" before these fields in order to assemble correctly.

The assembly language program is formulated as a sequence of statements of the above form, terminated optionally by an END statement. All statements following the END are ignored by the assembler.

### 3. FORMING THE OPERAND.

In order to completely describe the operation codes and pseudo operations, it is necessary to first present the form of the operand field, since it is used in nearly all statements. Expressions in the operand field consist of simple operands (labels, constants, and reserved words), combined in properly formed subexpressions by arithmetic and logical operators. The expression computation is carried out by the assembler as the assembly proceeds. Each expression must produce a 16-bit value during the assembly. Further, the number of significant digits in the result must not exceed the intended use. That is, if an expression is to be used in a byte move immediate instruction, then the most significant 8 bits of the expression must be zero. The restrictions on the expression significance is given with the individual instructions.

#### 3.1. Labels.

As discussed above, a label is an identifier which occurs on a particular statement. In general, the label is given a value determined by the type of statement which it precedes. If the label occurs on a statement which generates machine code or reserves memory space (e.g, a MOV instruction, or a DS pseudo operation), then the label is given the value of the program address which it labels. If the label precedes an EQU or SET, then the label is given the value which results from evaluating the operand field. Except for the SET statement, an identifier can label only one statement.

When a label appears in the operand field, its value is substituted by the assembler. This value can then be combined with other operands and operators to form the operand field for a particular instruction.

#### 3.2. Numeric Constants.

A numeric constant is a 16-bit value in one of several bases. The base, called the radix of the constant, is denoted by a trailing radix indicator. The radix indicators are

- B     binary constant (base 2)
- O     octal constant (base 8)



Gs,b,c

G,b

G,b,c

The first form starts execution of the program under test at the current value of the program counter in the current machine state, with no breakpoints set (the only way to regain control in DDT is through a RST 7 execution). The current program counter can be viewed by typing an X or XP command. The second form is similar to the first except that the program counter in the current machine state is set to address s before execution begins. The third form is the same as the second, except that program execution stops when address b is encountered (b must be in the area of the program under test). The instruction at location b is not executed when the breakpoint is encountered. The fourth form is identical to the third, except that two breakpoints are specified, one at b and the other at c. Encountering either breakpoint causes execution to stop, and both breakpoints are subsequently cleared. The last two forms take the program counter from the current machine state, and set one and two breakpoints, respectively.

Execution continues from the starting address in real-time to the next breakpoint. That is, there is no intervention between the starting address and the break address by DDT. Thus, if the program under test does not reach a breakpoint, control cannot return to DDT without executing a RST 7 instruction. Upon encountering a breakpoint, DDT stops execution and types

\*d

where d is the stop address. The machine state can be examined at this point using the X (Examine) command. The operator must specify breakpoints which differ from the program counter address at the beginning of the G command. Thus, if the current program counter is 1234H, then the commands

G,1234

and

G400,400

both produce an immediate breakpoint, without executing any instructions whatsoever.

5. The I (Input) Command. The I command allows the operator to insert a file name into the default file control block at 5CH (the file control block created by CP/M for transient programs is placed at this location; see the CP/M Interface Guide). The default FCB can be used by the program under test as if it had been passed by the CP/M Console Processor. Note that this file name is also used by DDT for reading additional HEX and COM files. The form of the I command is

Ifilename

or



of the instruction (e.g, MOV A,B), the value of the instruction (in this case MOV) is the bit pattern of the instruction with zeroes in the optional fields (e.g, MOV produces 40H).

When the symbol "\$" occurs in the operand field (not imbedded within identifiers and numeric constants) its value becomes the address of the next instruction to generate, not including the instruction contained within the current logical line.

### 3.4. String Constants.

String constants represent sequences of ASCII characters, and are represented by enclosing the characters within apostrophe symbols ('). All strings must be fully contained within the current physical line (thus allowing "!" symbols within strings), and must not exceed 64 characters in length. The apostrophe character itself can be included within a string by representing it as a double apostrophe (the two keystrokes ''), which becomes a single apostrophe when read by the assembler. In most cases, the string length is restricted to either one or two characters (the DB pseudo operation is an exception), in which case the string becomes an 8 or 16 bit value, respectively. Two character strings become a 16-bit constant, with the second character as the low order byte, and the first character as the high order byte.

The value of a character is its corresponding ASCII code. There is no case translation within strings, and thus both upper and lower case characters can be represented. Note however, that only graphic (printing) ASCII characters are allowed within strings. Valid strings are

```
'A'      'AB'      'ab'      'c'
'...'    'a'
'Walla Walla Wash.'
'She said "Hello" to me.'
'I said "Hello" to her.'
```

### 3.5. Arithmetic and Logical Operators.

The operands described above can be combined in normal algebraic notation using any combination of properly formed operands, operators, and parenthesized expressions. The operators recognized in the operand field are

a + b	unsigned arithmetic sum of a and b
a - b	unsigned arithmetic difference between a and b
+ b	unary plus (produces b)
- b	unary minus (identical to 0 - b)
a * b	unsigned magnitude multiplication of a and b
a / b	unsigned magnitude division of a by b
a MOD b	remainder after a / b
NOT b	logical inverse of b (all 0's become 1's, 1's become 0's), where b is considered a 16-bit value



assuming the tested program does not destroy the default area at 5CH. Further, any file specified with the filetype "COM" is assumed to contain machine code in pure binary form (created with the LOAD or SAVE command), and all others are assumed to contain machine code in Intel hex format (produced, for example, with the ASM command).

Recall that the command

```
DDT filename.filetype
```

which initiates the DDT program is equivalent to the commands

```
DDT
-Ifilename.filetype
-R
```

Whenever the R command is issued, DDT responds with either the error indicator "?" (file cannot be opened, or a checksum error occurred in a HEX file), or with a load message taking the form

```
NEXT PC
nnnn pppp
```

where nnnn is the next address following the loaded program, and pppp is the assumed program counter (100H for COM files, or taken from the last record if a HEX file is specified).

9. The S (Set) Command. The S command allows memory locations to be examined and optionally altered. The form of the command is

```
Ss
```

where s is the hexadecimal starting address for examination and alteration of memory. DDT responds with a numeric prompt, giving the memory location, along with the data currently held in the memory location. If the operator types a carriage return, then the data is not altered. If a byte value is typed, then the value is stored at the prompted address. In either case, DDT continues to prompt with successive addresses and values until either a period (.) is typed by the operator, or an invalid input value is detected.

10. The T (Trace) Command. The T command allows selective tracing of program execution for 1 to 65535 program steps. The forms are

```
T
Tn
```

In the first case, the CPU state is displayed, and the next program step is executed. The program terminates immediately, with the termination address



displayed as

\*hhhh

where hhhh is the next address to execute. The display address (used in the D command) is set to the value of H and L, and the list address (used in the L command) is set to hhhh. The CPU state at program termination can then be examined using the X command.

The second form of the T command is similar to the first, except that execution is traced for n steps (n is a hexadecimal value) before a program breakpoint is occurs. A breakpoint can be forced in the trace mode by typing a rubout character. The CPU state is displayed before each program step is taken in trace mode. The format of the display is the same as described in the X command.

Note that program tracing is discontinued at the interface to CP/M, and resumes after return from CP/M to the program under test. Thus, CP/M functions which access I/O devices, such as the diskette drive, run in real-time, avoiding I/O timing problems. Programs running in trace mode execute approximately 500 times slower than real time since DDT gets control after each user instruction is executed. Interrupt processing routines can be traced, but it must be noted that commands which use the breakpoint facility (G, T, and U) accomplish the break using a RST 7 instruction, which means that the tested program cannot use this interrupt location. Further, the trace mode always runs the tested program with interrupts enabled, which may cause problems if asynchronous interrupts are received during tracing.

Note also that the operator should use the rubout key to get control back to DDT during trace, rather than executing a RST 7, in order to ensure that the trace for the current instruction is completed before interruption.

11. The U (Untrace) Command. The U command is identical to the T command except that intermediate program steps are not displayed. The untrace mode allows from 1 to 65535 (0FFFFH) steps to be executed in monitored mode, and is used principally to retain control of an executing program while it reaches steady state conditions. All conditions of the T command apply to the U command.

12. The X (Examine) Command. The X command allows selective display and alteration of the current CPU state for the program under test. The forms are

X  
Xr

where r is one of the 8080 CPU registers

C	Carry Flag	(0/1)
Z	Zero Flag	(0/1)



M	Minus Flag	(0/1)
E	Even Parity Flag	(0/1)
I	Interdigit Carry	(0/1)
A	Accumulator	(0-FF)
B	BC register pair	(0-FFFF)
D	DE register pair	(0-FFFF)
H	HL register pair	(0-FFFF)
S	Stack Pointer	(0-FFFF)
P	Program Counter	(0-FFFF)

In the first case, the CPU register state is displayed in the format

CfzfMfEfIf A=bb B=dddd D=dddd H=dddd S=dddd P=dddd inst

where f is a 0 or 1 flag value, bb is a byte value, and dddd is a double byte quantity corresponding to the register pair. The "inst" field contains the disassembled instruction which occurs at the location addressed by the CPU state's program counter.

The second form allows display and optional alteration of register values, where r is one of the registers given above (C, Z, M, E, I, A, B, D, H, S, or P). In each case, the flag or register value is first displayed at the console. The DDT program then accepts input from the console. If a carriage return is typed, then the flag or register value is not altered. If a value in the proper range is typed, then the flag or register value is altered. Note that BC, DE, and HL are displayed as register pairs. Thus, the operator types the entire register pair when B, C, or the BC pair is altered.

### III. IMPLEMENTATION NOTES.

The organization of DDT allows certain non-essential portions to be overlayed in order to gain a larger transient program area for debugging large programs. The DDT program consists of two parts: the DDT nucleus and the assembler/disassembler module. The DDT nucleus is loaded over the Console Command Processor, and, although loaded with the DDT nucleus, the assembler/disassembler is overlayable unless used to assemble or disassemble.

In particular, the BDOS address at location 6H (address field of the JMP instruction at location 5H) is modified by DDT to address the base location of the DDT nucleus which, in turn, contains a JMP instruction to the BDOS. Thus, programs which use this address field to size memory see the logical end of memory at the base of the DDT nucleus rather than the base of the BDOS.

The assembler/disassembler module resides directly below the DDT nucleus in the transient program area. If the A, L, T, or X commands are used during the debugging process then the DDT program again alters the address field at 6H to include this module, thus further reducing the logical end of memory. If a program loads beyond the beginning of the assembler/disassembler module, the A and L commands are lost (their use produces a "?" in response), and the



trace and display (T and X) commands list the "inst" field of the display in hexadecimal, rather than as a decoded instruction.

#### IV. AN EXAMPLE.

The following example shows an edit, assemble, and debug for a simple program which reads a set of data values and determines the largest value in the set. The largest value is taken from the vector, and stored into "LARGE" at the termination of the program

```

ED SCAN.ASM,
*I, ↑-I ORG ↑-I 100H L-L START OF TRANSIENT AREA,
      MVI B,LEN LENGTH OF VECTOR TO SCAN,
      MVI C,0 LARGEST-EST VALUE SO FAR,
LOOP: P-0-0-L LXI H,VECT BASE OF VECTOR,
      MOV A,M GET VALUE,
      SUB C LARGER VALUE IN C?
      JNC NFOUND JUMP IF LARGER VALUE NOT FOUND,
      NEW LARGEST VALUE, STORE IT TO C,
      MOV C,A
NFOUND: INX H TO NEXT ELEMENT,
      DCR B MORE TO SCAN?
      JNZ LOOP FOR ANOTHER,

      END OF SCAN, STORE C,
      MOV A,C GET LARGEST VALUE,
      STA LARGE,
      JMP 0 REBOOT,

      TEST DATA
VECT: DB 2,0,4,3,5,6,1,5,
LEN: EQU $-VECT LENGTH,
LARGE: DS 1 LARGEST VALUE ON EXIT,
      END

*BOP,
      ORG 100H START OF TRANSIENT AREA
      MVI B,LEN LENGTH OF VECTOR TO SCAN
      MVI C,0 LARGEST VALUE SO FAR
      LXI H,VECT BASE OF VECTOR
LOOP: MOV A,M GET VALUE
      SUB C LARGER VALUE IN C?
      JNC NFOUND JUMP IF LARGER VALUE NOT FOUND
      NEW LARGEST VALUE, STORE IT TO C
      MOV C,A
NFOUND: INX H TO NEXT ELEMENT
      DCR B MORE TO SCAN?
      JNZ LOOP FOR ANOTHER

```

*Handwritten notes:*

- tab character* (pointing to ↑-I)
- tabout* (pointing to L-L)
- tabout echo* (pointing to ↑-I)
- Reboot deletes characters* (pointing to ↑-I)

*Annotations:*

- Create Source Program - underlined characters typed by programmer.*
- "↵" represents carriage return.*



```

;      END OF SCAN, STORE C
      MOV      A,C      ;GET LARGEST VALUE
      STA      LARGE
      JMP      0        ;REBOOT

;
;      TEST DATA
VECT:  DB      2,0,4,3,5,6,1,5
LEN    EQU      $-VECT  ;LENGTH
LARGE: DS      1        ;LARGEST VALUE ON EXIT
      END

```

\*E, ← End of Edit

ASM SCAN, Start Assembler  
 CP/M ASSEMBLER - VER 1.0

0122  
 002H USE FACTOR  
 END OF ASSEMBLY

Assembly Complete - Look at Program Listing

TYPE SCAN.PRN,

Code Address	Machine Code	Source Program
0100		ORG 100H ; START OF TRANSIENT AREA
0100 0608		MVI B,LEN ; LENGTH OF VECTOR TO SCAN
0102 0E00		MVI C,0 ; LARGEST VALUE SO FAR
0104 211901		LXI H,VECT ; BASE OF VECTOR
0107 7E	LOOP:	MOV A,M ; GET VALUE
0108 91		SUB C ; LARGER VALUE IN C?
0109 D20D01		JNC NFOUND ; JUMP IF LARGER VALUE NOT FOUND
		NEW LARGEST VALUE, STORE IT TO C
010C 4F		MOV C,A
010D 23	NFOUND:	INX H ; TO NEXT ELEMENT
010E 05		DCR B ; MORE TO SCAN?
010F C20701		JNZ LOOP ; FOR ANOTHER
		END OF SCAN, STORE C
0112 79		MOV A,C ; GET LARGEST VALUE
0113 322101		STA LARGE
0116 C30000		JMP 0 ; REBOOT
		TEST DATA
0119 0200040305	VECT:	DB 2,0,4,3,5,6,1,5
0008 =	LEN	EQU \$-VECT ; LENGTH
0121 Value of	LARGE:	DS 1 ; LARGEST VALUE ON EXIT
0122 Equate		END

A>

DDT SCAN.HEX,

Start Debugger using hex format machine code

16K DDT VER 1.0

NEXT PC

0121 0000

-X, last load address + 1

0020M0E010 A=00 B=0000 D=0000 H=0000 S=0100 P=0000 OUT 7F

-XP,

Examine registers before debug run

P=0000 100,

Change PC to 100

-X, Look at registers again

PC changed.

0020M0E010 A=00 B=0000 D=0000 H=0000 S=0100 P=0100 MVI B,08

-L100,

Next instruction  
to execute at PC=100

0100 MVI B,08  
0102 MVI C,00  
0104 LXI H,0119  
0107 MOV A,M  
0108 SUB C  
0109 JNC 010D  
010C MOV C,A  
010D INX H  
010E DCR B  
010F JNZ 0107  
0112 MOV A,C

Disassembled Machine  
Code at 100H  
(See Source Listing  
for comparison)

-L,

0113 STA 0121  
0116 JMP 0000  
0119 STAX B  
011A NOP  
011B INR B  
011C INX B  
011D DCR B  
011E MVI B,01  
0120 DCR B  
0121 LXI D,2200  
0124 LXI H,0200

A little more  
machine code  
(note that program  
ends at location 116  
with a JMP to 0000)

-A116, enter inline assembly mode to change the JMP to 0000 into a RST 7, which  
will cause the program under test to return to DDT if 116H  
is ever executed.

0116 RST 7,

0117, (single carriage return stops assemble mode)

-L113, List Code at 113H to check that RST 7 was properly inserted

0113 STA 0121

0116 RST 07

In Place of JMP



```

0117 NOP
0118 NOP
0119 STAX B
011A NOP
011B INR B
011C INX B

```

-X, Look at registers

```
C0Z0M0E010 A=00 B=0000 D=0000 H=0000 S=0100 P=0100 MVI B,08
```

-I, Execute Program for one step. initial CPU state, before is executed

```
C0Z0M0E010 A=00 B=0000 D=0000 H=0000 S=0100 P=0100 MVI B,08*0102
```

-I, Trace one step again (note 08H in B) automatic breakpoint

```
C0Z0M0E010 A=00 B=0800 D=0000 H=0000 C=0100 P=0102 MVI C,00*0104
```

-I, Trace again (Register C is cleared)

```
C0Z0M0E010 A=00 B=0800 D=0000 H=0000 S=0100 P=0104 LXI H,0119*0107
```

-T3, Trace three steps

```
C0Z0M0E010 A=00 B=0800 D=0000 H=0119 S=0100 P=0107 MOV A,M
```

```
C0Z0M0E010 A=02 B=0800 D=0000 H=0119 S=0100 P=0108 SUB C
```

```
C0Z0M0E011 A=02 B=0800 D=0000 H=0119 S=0100 P=0109 JNC 010D*010D
```

-D119, Display memory starting at 119H.

Address	02	00	04	03	05	06	01	Program data
0119	02	00	04	03	05	06	01	
0120	05	11	00	22	21	00	02	7E EB 77 13 23 EB 0B 78 B1
0130	C2	27	01	C3	03	29	00	00 00 00 00 00 00 00 00
0140	00	00	00	00	00	00	00	00 00 00 00 00 00 00 00
0150	00	00	00	00	00	00	00	00 00 00 00 00 00 00 00
0160	00	00	00	00	00	00	00	00 00 00 00 00 00 00 00
0170	00	00	00	00	00	00	00	00 00 00 00 00 00 00 00
0180	00	00	00	00	00	00	00	00 00 00 00 00 00 00 00
0190	00	00	00	00	00	00	00	00 00 00 00 00 00 00 00
01A0	00	00	00	00	00	00	00	00 00 00 00 00 00 00 00
01B0	00	00	00	00	00	00	00	00 00 00 00 00 00 00 00
01C0	00	00	00	00	00	00	00	00 00 00 00 00 00 00 00

Automatic breakpoint at 10DH

Lowercase x

Data is displayed in ASCII with a "0" in the position of non-graphic characters

-X, Current CPU state

```
C0Z0M0E011 A=02 B=0800 D=0000 H=0119 S=0100 P=010D INX H
```

-T5, Trace 5 steps from current CPU state

```
C0Z0M0E011 A=02 B=0800 D=0000 H=0119 S=0100 P=010D INX H
```

```
C0Z0M0E011 A=02 B=0800 D=0000 H=011A S=0100 P=010E DCR B
```

```
C0Z0M0E011 A=02 B=0700 D=0000 H=011A S=0100 P=010F JNZ 0107
```

```
C0Z0M0E011 A=02 B=0700 D=0000 H=011A S=0100 P=0107 MOV A,M
```

```
C0Z0M0E011 A=00 B=0700 D=0000 H=011A S=0100 P=0108 SUB C*0109
```

-U5, Trace without listing intermediate states

```
C0Z1M0E111 A=00 B=0700 D=0000 H=011A S=0100 P=0109 JNC 010D*0108
```

-X, CPU state at end of U5

```
C0Z0M0E111 A=04 B=0600 D=0000 H=011B S=0100 P=0108 SUB C
```



$-\frac{x}{2}$  CPU state at end of Program

-X<sub>P</sub>: examine and change Program Counter

$$P = 0.116 \frac{100}{\text{}} ,$$
$$-x_2$$

-T10 Trace 10 (hexadecimal) steps

C0Z1M0E111	A=00	B=0000	D=0000	H=0121	S=0100	P=0100	MVI	C,00
C0Z1M0E111	A=00	B=0000	D=0000	H=0121	S=0100	P=0102	MVI	C,00
C0Z1M0E111	A=00	B=0000	D=0000	H=0121	S=0100	P=0104	LXI	H,0119
C0Z1M0E111	A=00	B=0000	D=0000	H=0119	S=0100	P=0107	MOV	A,M
C0Z1M0E111	A=02	B=0000	D=0000	H=0119	S=0100	P=0108	SUB	C
C0Z0M0E011	A=02	B=0000	D=0000	H=0119	S=0100	P=0109	JNC	010D
C0Z0M0E011	A=02	B=0000	D=0000	H=0119	S=0100	P=010D	INX	H
C0Z0M0E011	A=02	B=0000	D=0000	H=011A	S=0100	P=010E	DCR	B
C0Z0M0E011	A=02	B=0700	D=0000	H=011A	S=0100	P=010F	JNZ	0107
C0Z0M0E011	A=02	B=0700	D=0000	H=011A	S=0100	P=0107	MOV	A,M
C0Z0M0E011	A=00	B=0700	D=0000	H=011A	S=0100	P=0108	SUB	C
C0Z1M0E111	A=00	B=0700	D=0000	H=011A	S=0100	P=0109	JNC	010D
C0Z1M0E111	A=00	B=0700	D=0000	H=011A	S=0100	P=010D	INX	H
C0Z1M0E111	A=00	B=0700	D=0000	H=011B	S=0100	P=010E	DCR	B
C0Z0M0E111	A=00	B=0600	D=0000	H=011B	S=0100	P=010F	JNZ	0107
C0Z0M0E111	A=00	B=0600	D=0000	H=011B	S=0100	P=0107	MOV	A,M*0108

-A109, Insert a "hot patch" into  
0109 JC 10D, the machine code  
010C, to change the  
JWC to JC

-GG, Stop DDT so that a version of the patched program can be saved

Program should have moved the value from A into C since  $A > C$ . Since this code was not executed, it appears that the JNC should have been a JC instruction

SAVE 1 SCAN.COM → Program resides on first page, so save 1 page.

A>DDT SCAN.COM, Restart DDT with the saved memory image to continue testing

0200 0100

-L100, List some code

```
0100 MVI B,08
0102 MVI C,00
0104 LXI H,0119
0107 MOV A,M
0108 SUB C
0109 JC 010D
```

Previous Patch is Present in X.COM



```

010C MOV C,A
010D INX H
010E DCR B
010F JNZ 0107
0112 MOV A,C

```

-X,

P=0100,

-T10, Trace to see how patched version operates Data is moved from A to C

```

C0Z0M0E010 A=00 B=0000 D=0000 H=0000 S=0100 P=0100 MVI B,00
C0Z0M0E010 A=00 B=0000 D=0000 H=0000 S=0100 P=0102 MVI C,00
C0Z0M0E010 A=00 B=0000 D=0000 H=0000 S=0100 P=0104 LXI H,0119
C0Z0M0E010 A=00 B=0000 D=0000 H=0119 S=0100 P=0107 MOV A,M
C0Z0M0E010 A=02 B=0000 D=0000 H=0119 S=0100 P=0108 SUB C
C0Z0M0E011 A=02 B=0000 D=0000 H=0119 S=0100 P=0109 JC 010D
C0Z0M0E011 A=02 B=0000 D=0000 H=0119 S=0100 P=010C MOV C,A
C0Z0M0E011 A=02 B=0002 D=0000 H=0119 S=0100 P=010D INX H
C0Z0M0E011 A=02 B=0002 D=0000 H=011A S=0100 P=010E DCR B
C0Z0M0E011 A=02 B=0702 D=0000 H=011A S=0100 P=010F JNZ 0107
C0Z0M0E011 A=02 B=0702 D=0000 H=011A S=0100 P=0107 MOV A,M
C0Z0M0E011 A=00 B=0702 D=0000 H=011A S=0100 P=0108 SUB C
C1Z0M1E010 A=FE B=0702 D=0000 H=011A S=0100 P=0109 JC 010D
C1Z0M1E010 A=FE B=0702 D=0000 H=011A S=0100 P=010D INX H
C1Z0M1E010 A=FE B=0702 D=0000 H=011B S=0100 P=010E DCR B
C1Z0M0E111 A=FE B=0602 D=0000 H=011B S=0100 P=010F JNZ 0107*0107

```

-X,

breakpoint after 16 steps

C1Z0M0E111 A=FE B=0602 D=0000 H=011B S=0100 P=0107 MOV A,M

-G,108, Run from current PC and breakpoint at 108H

\*0108

-X,

next data item

C1Z0M0E111 A=04 B=0602 D=0000 H=011B S=0100 P=0108 SUB C

-T,

Single Step for a few cycles

C1Z0M0E111 A=04 B=0602 D=0000 H=011B S=0100 P=0108 SUB C\*0109

-T,

C0Z0M0E011 A=02 B=0602 D=0000 H=011B S=0100 P=0109 JC 010D\*010C

-X,

C0Z0M0E011 A=02 B=0602 D=0000 H=011B S=0100 P=010C MOV C,A

-G, Run to completion

\*0116

-X,

C0Z1M0E111 A=03 B=0003 D=0000 H=0121 S=0100 P=0116 RST 07

-S121, look at the value of "LARGE"

0121 03, Wrong Value!



0122 00,

0123 22,

0124 21,

0125 00,

0126 02,

0127 7E,

End of the S Command

-L100,

0100 MVI B,08  
0102 MVI C,00  
0104 LXI H,0119  
0107 MOV A,M  
0108 SUB C  
0109 JC 010D  
010C MOV C,A  
010D INX H  
010E DCR B  
010F JNZ 0107  
0112 MOV A,C

Review the Code

-L,

0113 STA 0121  
0116 RST 07  
0117 NOP  
0118 NOP  
0119 STAX B  
011A NOP  
011B INR B  
011C INX B  
011D DCR B  
011E MVI B,01  
0120 DCR B

-XP,

P=0116 100, Reset the PC

-I, Single step, and watch data values

0021M0E111 A=03 B=0003 D=0000 H=0121 S=0100 P=0100 MVI B,08\*0102

-I,

0021M0E111 A=03 B=0803 D=0000 H=0121 S=0100 P=0102 MVI C,00\*0104

-I,

0021M0E111 A=03 B=0800 D=0000 H=0121 S=0100 P=0104 LXI H,0119\*0107

-I,

0021M0E111 A=03 B=0800 D=0000 H=0119 S=0100 P=0107 MOV A,M\*0108

-I,

first data item brought to A

C0Z1M0E111 A=02 B=0800 D=0000 H=0119 S=0100 P=0108 SUB C\*0109

-I,

C0Z0M0E011 A=02 B=0800 D=0000 H=0119 S=0100 P=0109 JC 010D\*010C

-I,

C0Z0M0E011 A=02 B=0800 D=0000 H=0119 S=0100 P=010C MOV C,A\*010D

-I,

first data item moved to C correctly

C0Z0M0E011 A=02 B=0802 D=0000 H=0119 S=0100 P=010D INX H\*010E

-I,

C0Z0M0E011 A=02 B=0802 D=0000 H=011A S=0100 P=010E DCR B\*010F

-I,

C0Z0M0E011 A=02 B=0702 D=0000 H=011A S=0100 P=010F JNZ 0107\*0107

-I,

C0Z0M0E011 A=02 B=0702 D=0000 H=011A S=0100 P=0107 MOV A,M\*0108

-I,

second data item brought to A

C0Z0M0E011 A=00 B=0702 D=0000 H=011A S=0100 P=0108 SUB C\*0109

-I,

subtract destroys data value which was loaded!!!

C1Z0M1E010 A=FE B=0702 D=0000 H=011A S=0100 P=0109 JC 010D\*010D

-I,

C1Z0M1E010 A=FE B=0702 D=0000 H=011A S=0100 P=010D INX H\*010E

-L100,

0100 MVI B,08  
0102 MVI C,00  
0104 LXI H,0119  
0107 MOV A,M  
0108 SUB C  
0109 JC 010D  
010C MOV C,A  
010D INX H  
010E DCR B  
010F JNZ 0107  
0112 MOV A,C

-A108,

0108 CMP C, hot patch at 108H changes sub to CMP

0109,

-G0, stop DDT for SAVE

← This should have been a CMP so that register A would not be destroyed.



SAVE 1 SCAN.COM,

Save memory image

A>DDT SCAN.COM,

Restart DDT

16K DDT VER 1.0

NEXT PC

0200 0100

-XP,

P=0100,

-L116,

0116 RST 07

0117 NOP

0118 NOP

0119 STAX B

011A NOP

- (rubout)

} Look at code to see if it was properly loaded  
(long timeout aborted with rubout)

-G.116, Run from look to completion

\*0116

-XC, Look at Carry (accidental typo)

C1,

-X, Look at CPU state

C1Z1M0E111 A=06 B=0006 D=0000 H=0121 S=0100 P=0116 RST 07

-S121, Look at "Large" - it appears to be correct.

0121 06,

0122 00,

0123 22 .,

-G0, stop DDT

ED SCAN.ASM,

Re-edit the source program, and make both changes

\*NSUB

\*0LT,

\*SSUBZCMPZ0LT,

\*,

\*SNCZCZ0LT,

\*E,

SUB

CMF

JNC

JC

C

C

NFOUND

NFOUND

LARGER VALUE IN C?

LARGER VALUE IN C?

JUMP IF LARGER VALUE NOT FOUND

JUMP IF LARGER VALUE NOT FOUND

ASM SCAN.AAZ, Re-assemble, selecting source from disk A  
 CP/M ASSEMBLER - VER 1.0 hex to disk A  
 Print to Z (selects no print file)

0122  
 002H USE FACTOR  
 END OF ASSEMBLY

DDT SCAN.HEX, Re-run debugger to check changes

16K DDT VER 1.0  
 NEXT PC  
 0121 0000  
 -L116,

0116 JMP 0000 check to ensure end is still at 116H  
 0119 STAX B  
 011A NOP  
 011B INR B  
 - (rubout)

-G100,116, Go from beginning with breakpoint at end

\*0116 breakpoint reached

-D121, Look at "LARGE" correct value computed

0121	06	00	22	21	00	02	7E	EB	77	13	23	EB	00	78	B1	...	"I...W...X"
0130	C2	27	01	C3	03	29	00	00	00	00	00	00	00	00	00	00	...
0140	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	...

- (rubout) aborts long typeout

-GO  
 -> stop DDT, debug session complete



1. The first part of the paper is devoted to a discussion of the general principles of the theory of the structure of the atom.

2. The second part of the paper is devoted to a discussion of the general principles of the theory of the structure of the atom.

3. The third part of the paper is devoted to a discussion of the general principles of the theory of the structure of the atom.

4. The fourth part of the paper is devoted to a discussion of the general principles of the theory of the structure of the atom.

5. The fifth part of the paper is devoted to a discussion of the general principles of the theory of the structure of the atom.

6. The sixth part of the paper is devoted to a discussion of the general principles of the theory of the structure of the atom.

7. The seventh part of the paper is devoted to a discussion of the general principles of the theory of the structure of the atom.

8. The eighth part of the paper is devoted to a discussion of the general principles of the theory of the structure of the atom.

9. The ninth part of the paper is devoted to a discussion of the general principles of the theory of the structure of the atom.





